

HyBF: A Hybrid Branch Fusion Strategy for Code Size Reduction

R. Rocha, C. Saumya, K. Sundararajah, P. Petoumenos, M. Kulkarni, M. O'Boyle



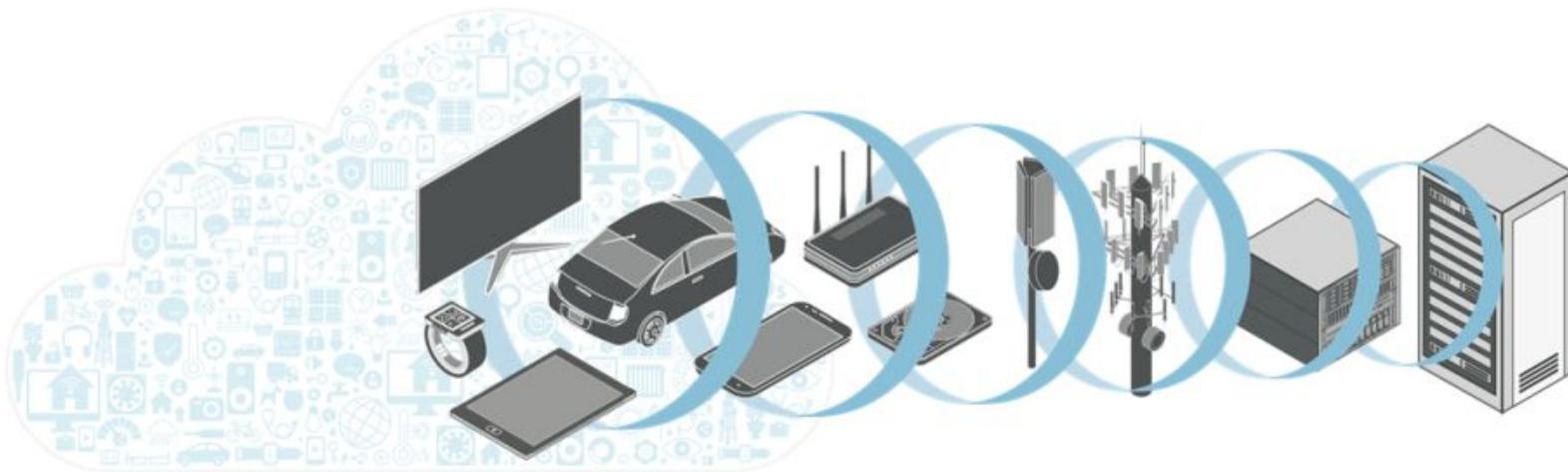
HyBF: A Hybrid Branch Fusion Strategy for Code Size Reduction

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Code Size Matters

Code size is important in many domains -- Large is relative to the constraints.

Compilers need powerful optimizations for code size.



What is Branch Fusion?

```
if (atomic_dec_test(&rcus.bcpucount)) {  
    rcu_btrace(TPS("LastCB"), -1, rcus.bseq);  
    complete(&rcus.bcompletion);  
} else {  
    rcu_btrace(TPS("CB"), -1, rcus.bseq);  
}
```

What is Branch Fusion?

Identify similarities between the *then* and *else* paths in conditional branches.

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What is Branch Fusion?

Merge similar code from the *then* and *else* paths of the branch.

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```
if (atomic_dec_test(&rcus.bcpucount)) {    cond = atomic_dec_test(&rcus.bcpucount)
    rcu_btrace(TPS("LastCB"), -1, rcus.bseq);    mval = cond ? "LastCB" : "CB";
    complete(&rcus.bcompletion);    rcu_btrace(TPS(mval), -1, rcus.bseq);
} else {
    rcu_btrace(TPS("CB"), -1, rcus.bseq);
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} else {    if (cond) {
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}
}
```

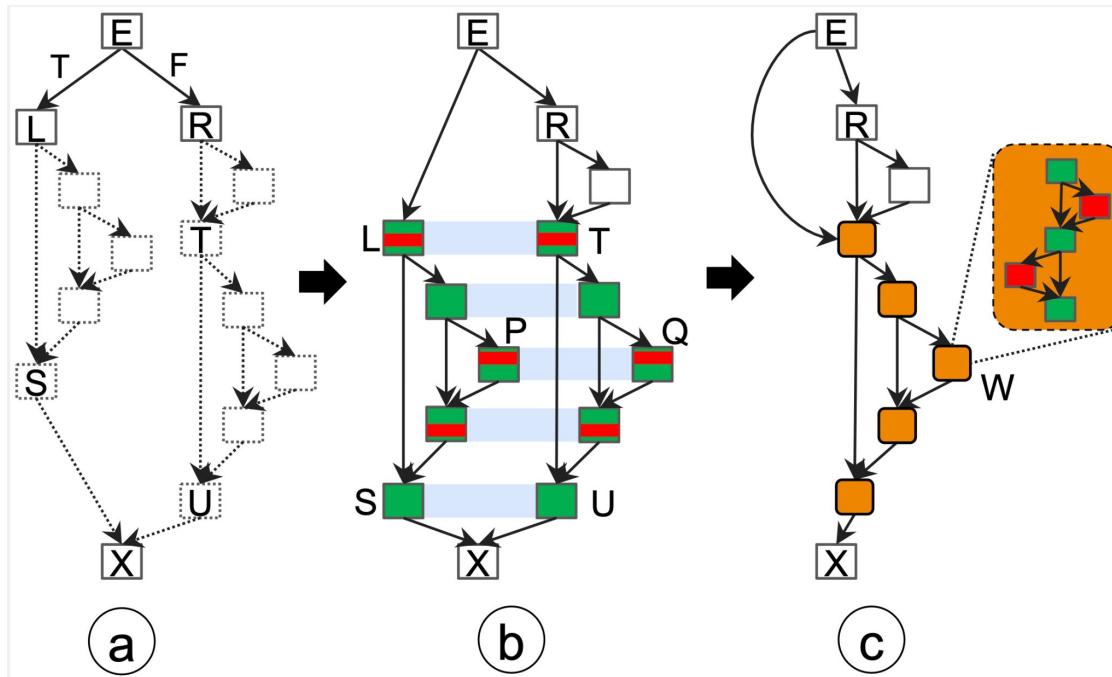
Main Insight: Branch Fusion Reduces Code Size

```
if (atomic_dec_test(&rcus.bcpucount)) {    cond = atomic_dec_test(&rcus.bcpucount)
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} else {                                if (cond) {
    rcu_btrace(TPS("CB"), -1, rcus.bseq);      complete(&rcus.bcompletion);
}                                            }
```

18 bytes (11%) Reduction

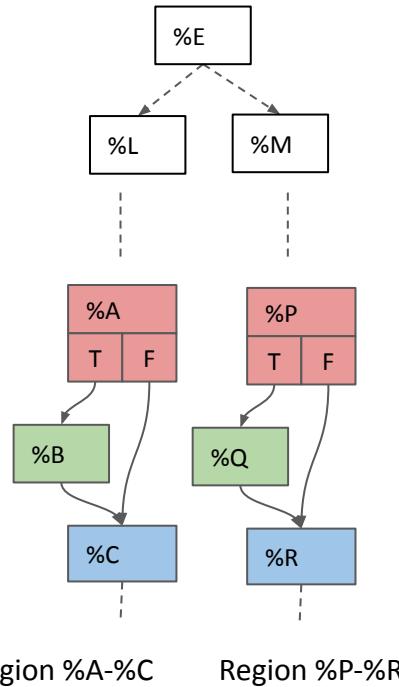
CFM-CS

Improved DARM to work on complex SESE regions in generic (non-GPU) code.



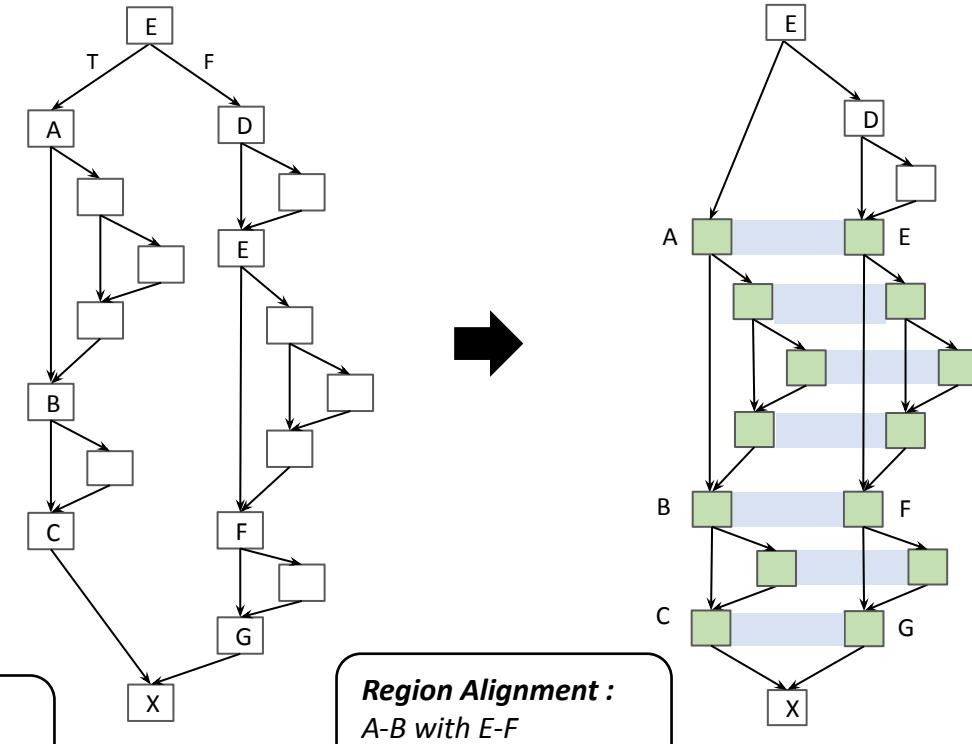
Meldable Regions

- Two SESE regions can be *melded* if,
 - Dominated by a conditional branch
 - No path exists that goes through both the SESE regions
 - Entry blocks of the regions must post-dominate either the left or right successor of the conditional branch
 - They are *isomorphic* (structural similarity)

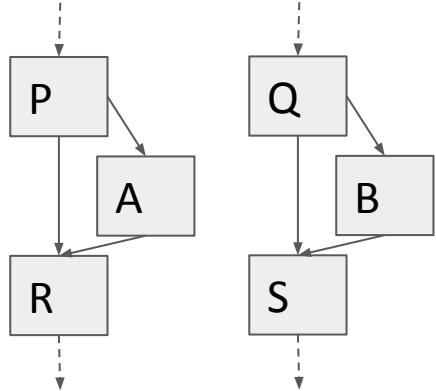


Region Alignment

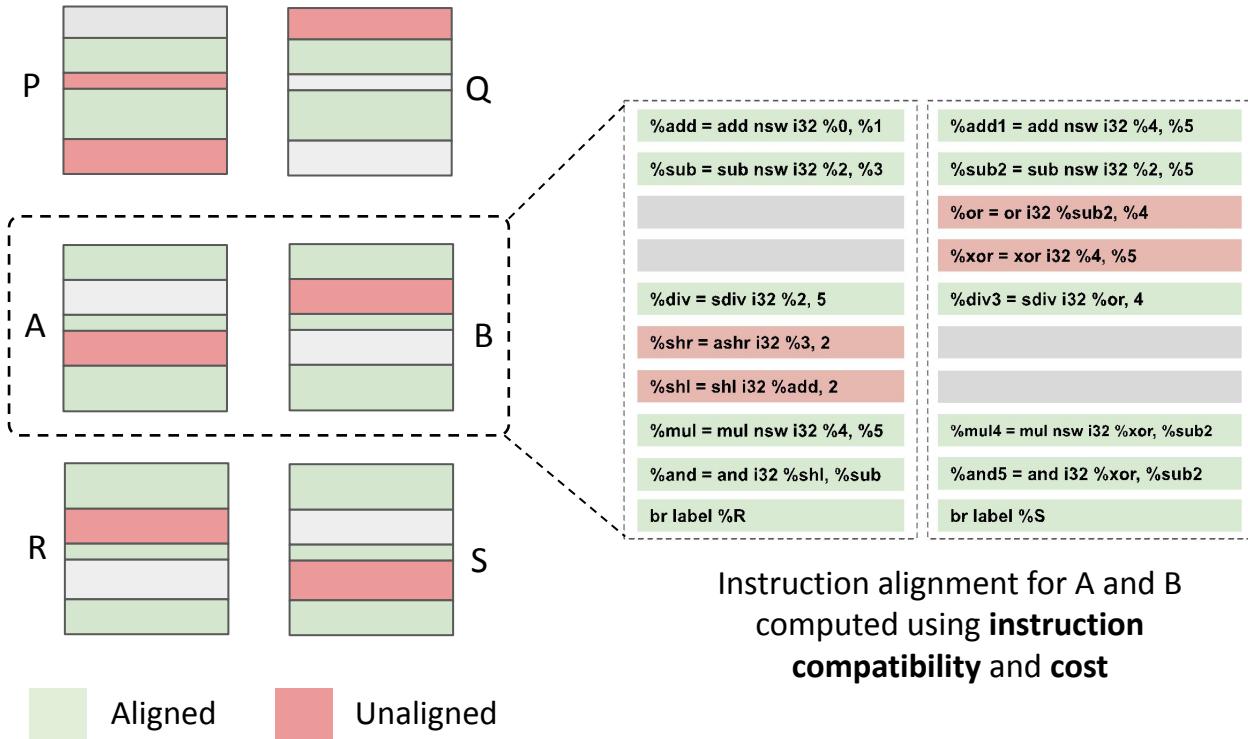
- Multiple isomorphic regions in if and else paths?
- Regions are aligned based on **Melding Profitability**
- **Melding Profitability** : metric that measures the similarity of two regions base on instruction frequencies



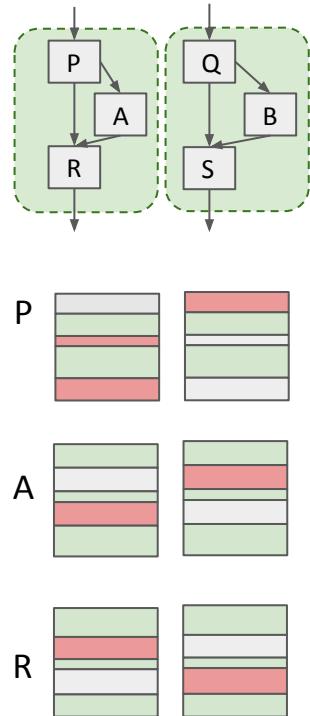
Instruction Alignment



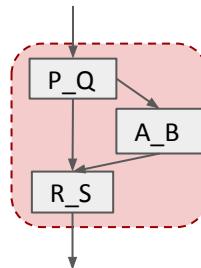
Aligned region pair



Code Generation



Generated melded control-flow



Generated melded instructions

Aligned instruction pair

%add = add nsw i32 %0, %1

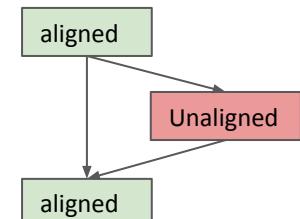
%add1 = add nsw i32 %4, %5

Generated code

%sel1 = select i1 %cmp, i32 %0, i32 %4

%sel2 = select i1 %cmp, i32 %1, i32 %5

%6 = add nsw i32 %sel1, %sel2



Unaligned instructions are executed conditionally

CFM-CS: Deep Dive Talk



Merging Similar Control-Flow Regions in LLVM for Performance and Code Size Benefits

Charitha Saumya (Presenter), Kirshanthan Sundararajah, Milind Kulkarni



LLVM Developers' Meeting 2022

November 7th – 10th 2022

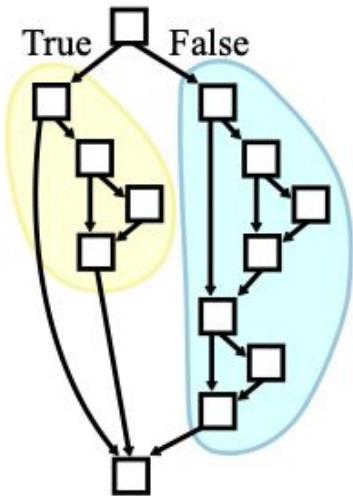
Merging Similar Control-Flow Regions in LLVM for
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Charitha Saumya

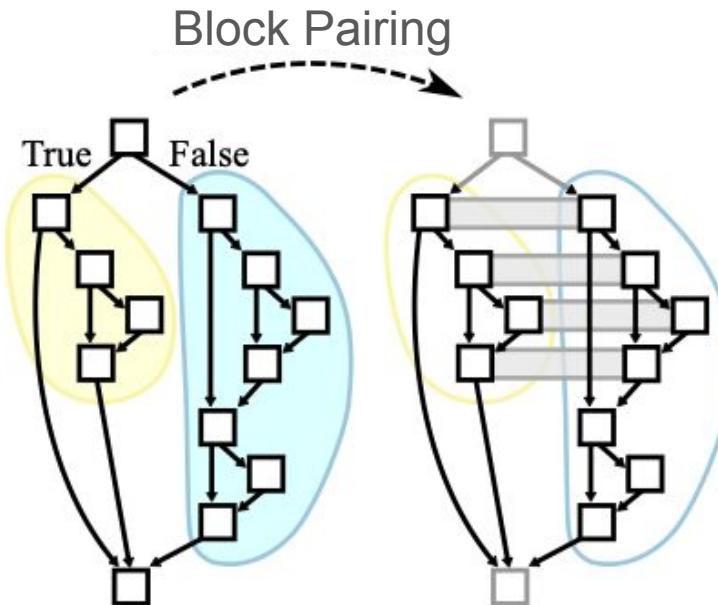
<https://www.youtube.com/watch?v=iGbdcltU0F8>

2022 LLVM Dev Mtg: Merging Similar Control-Flow Regions

Branch Fusion on SEME Regions

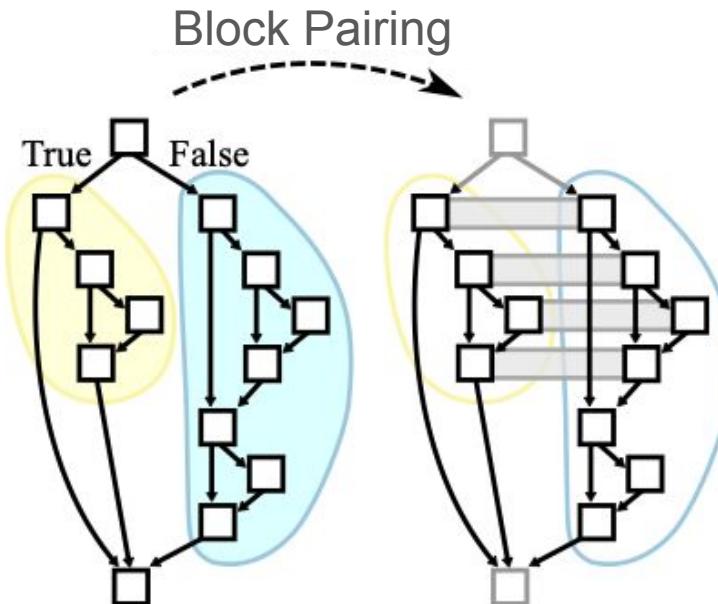


Branch Fusion on SEME Regions



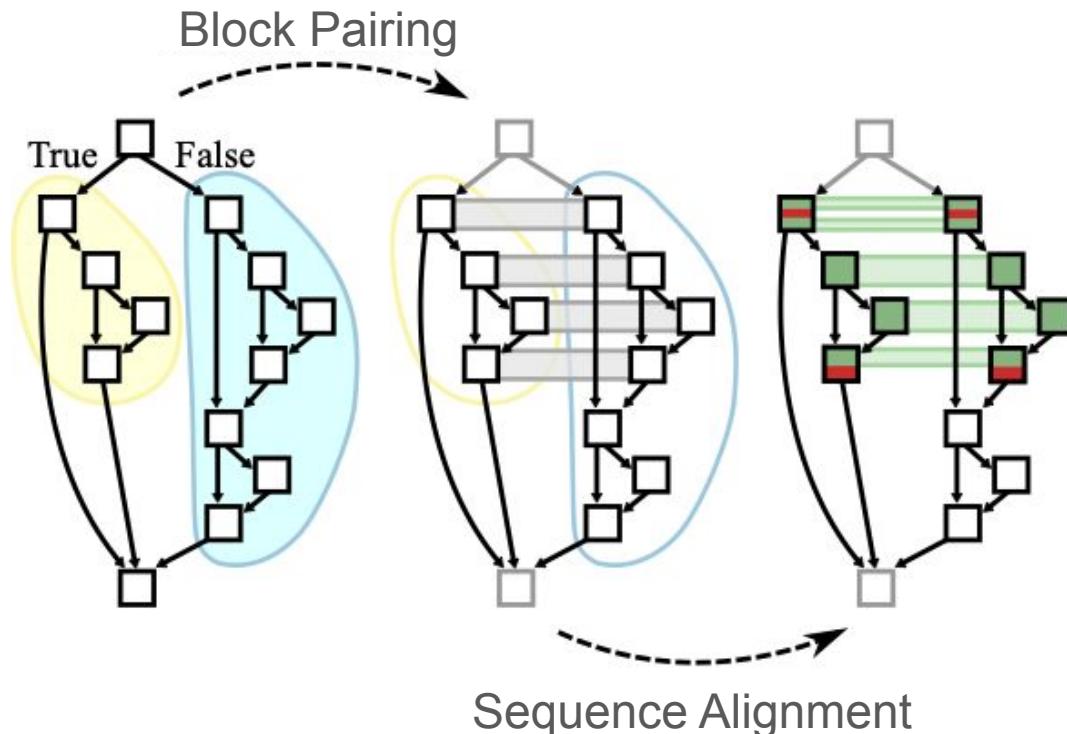
Blocks are paired independently, based only on their similarity.

Branch Fusion on SEME Regions

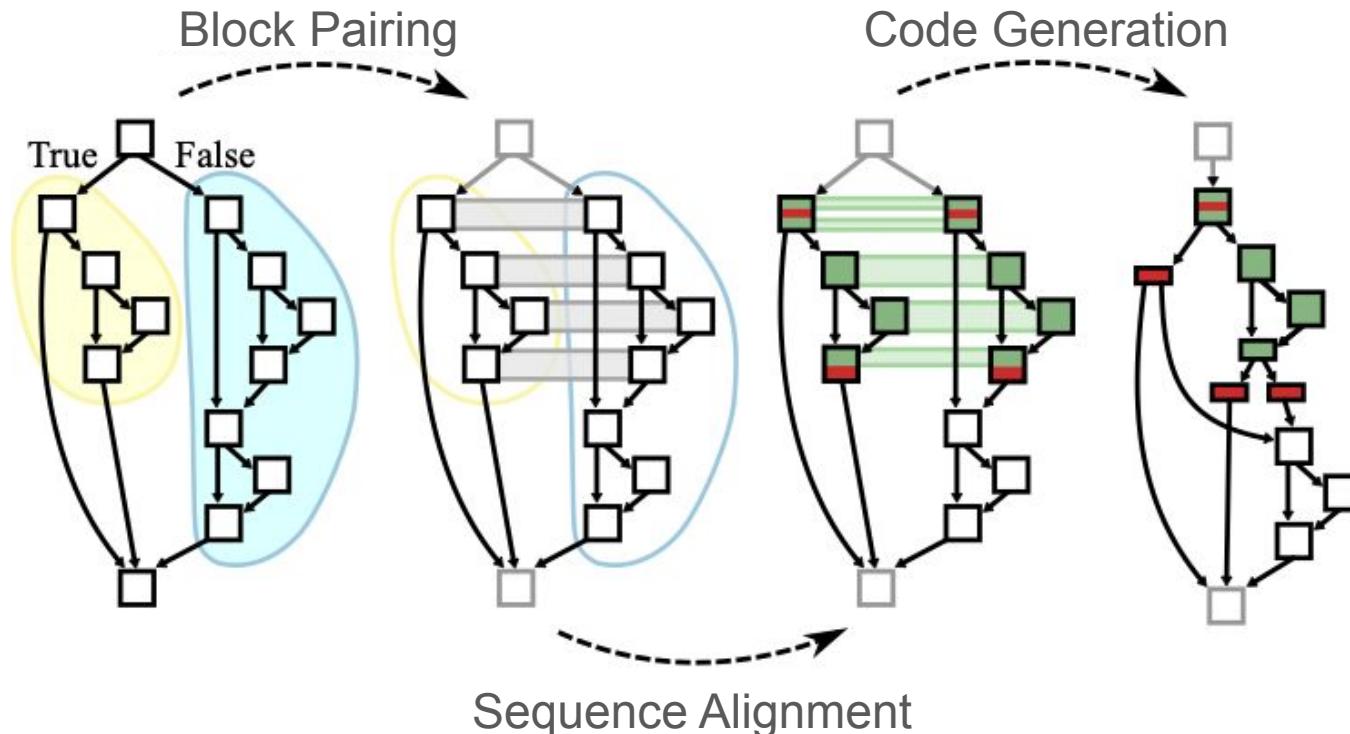


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Branch Fusion on SEME Regions



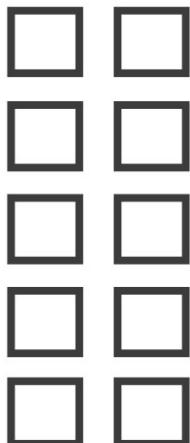
Branch Fusion on SEME Regions



Pairing Similar Blocks

Given two SEME regions

Region 1



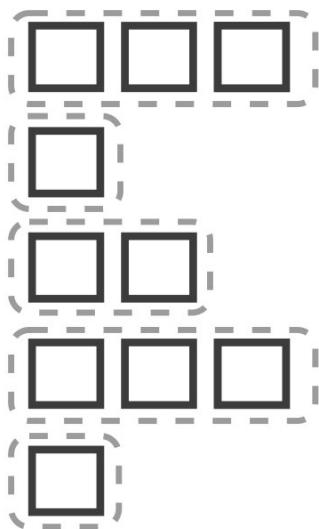
Region 2



Pairing Similar Blocks

Group blocks by their size

Region 1



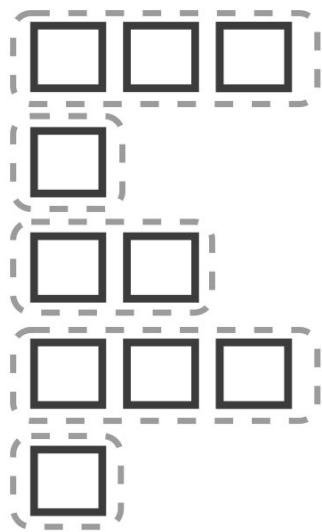
Region 2



Pairing Similar Blocks

Minimize fingerprint distance

Region 1



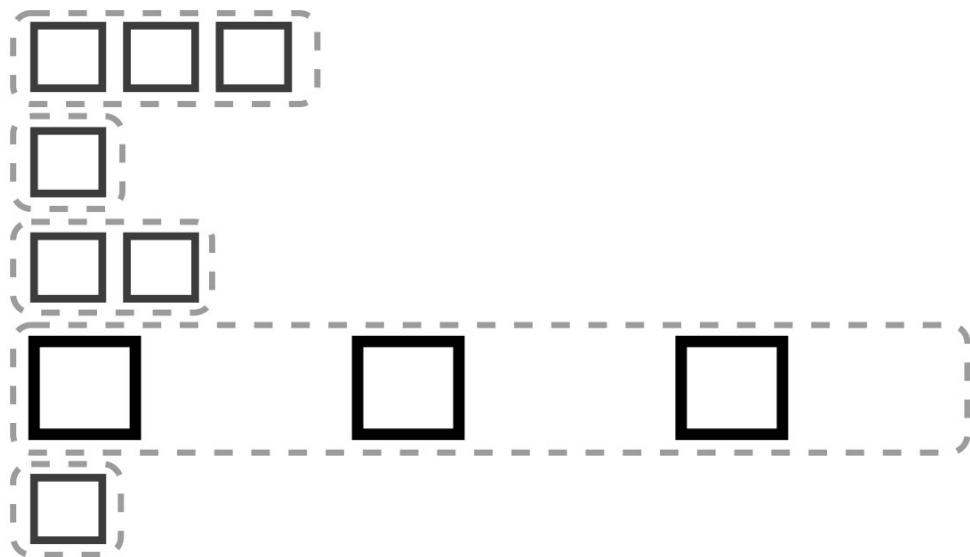
Region 2



Pairing Similar Blocks

Minimize fingerprint distance

Region 1



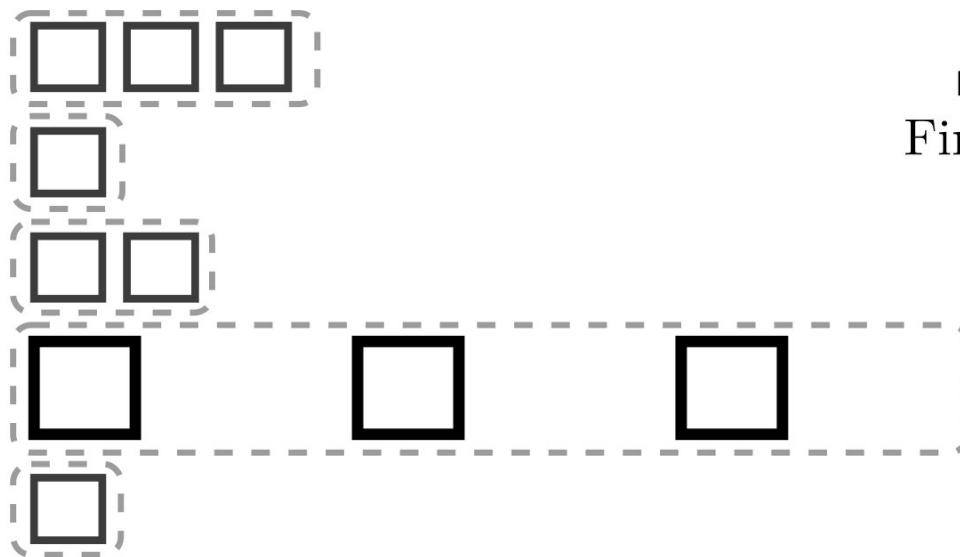
Region 2



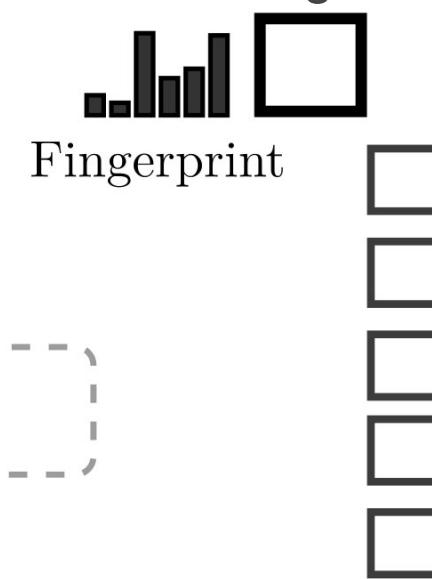
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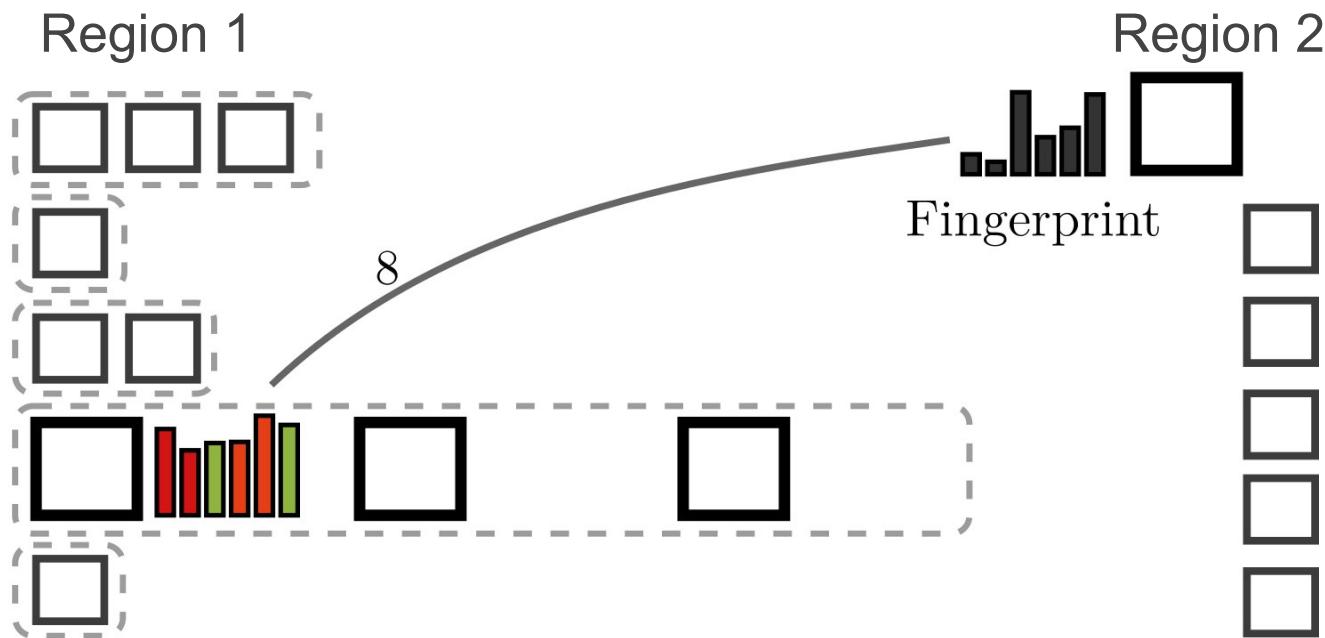


Region 2



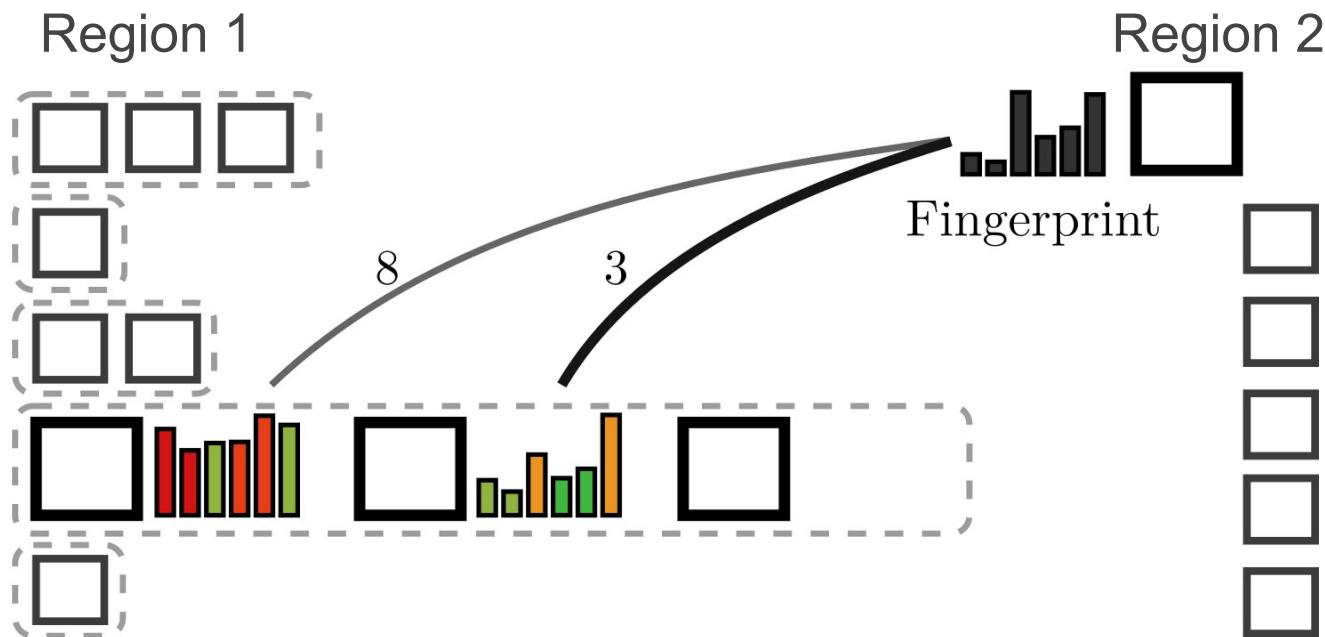
Pairing Similar Blocks

Minimize fingerprint distance



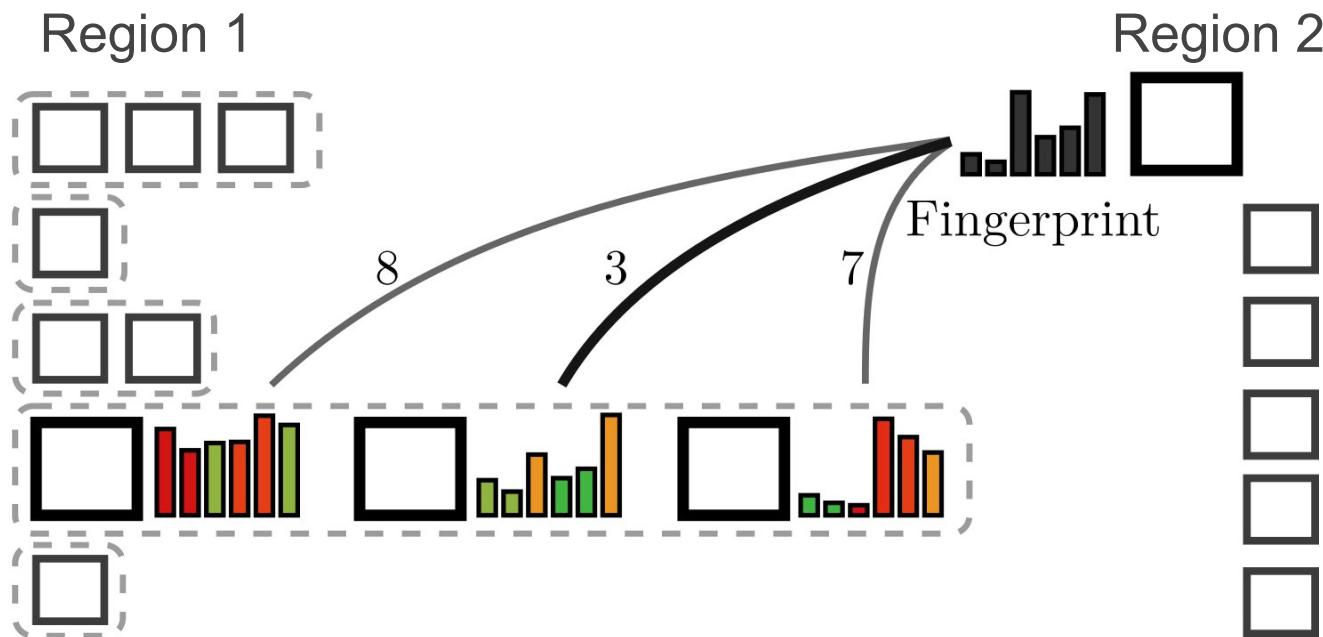
Pairing Similar Blocks

Minimize fingerprint distance



Pairing Similar Blocks

Minimize fingerprint distance



Pairwise Alignment of Paired Blocks

Match only corresponding instructions

%sw.bb	%entry
%v1 = gep %this, 0, 5	%x1 = alloca
%v2 = bitcast %v1	%x2 = gep %this, 0, 1
%v3 = load %v2	%x3 = load %x2
%v4 = icmp eq %v3, 0	%x4 = icmp eq %x3, 73
br %v4, L _{b3} , L _{b2}	br %x4, L _{b3} , L _{b2}

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Pairwise Alignment of Paired Blocks

Estimate size of merged block

%sw.bb	%entry	0
%v1 = gep %this, 0, 5	%x1 = alloca	
%v2 = bitcast %v1	%x2 = gep %this, 0, 1	
%v3 = load %v2	%x3 = load %x2	
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Pairwise Alignment of Paired Blocks

Estimate size of merged block

%sw.bb	%entry	0
%v1 = gep %this, 0, 5	%x1 = alloca	2
%v2 = bitcast %v1	%x2 = gep %this, 0, 1	
%v3 = load %v2	%x3 = load %x2	
%v4 = icmp eq %v3, 0	%x4 = icmp eq %x3, 73	
br %v4, L _{b3} , L _{b2}	br %x4, L _{b3} , L _{b2}	

Pairwise Alignment of Paired Blocks

Estimate size of merged block

%sw.bb	%entry	0) +1
%v1 = gep %this, 0, 5	%x1 = alloca	2	
%v2 = bitcast %v1	%x2 = gep %this, 0, 1		
%v3 = load %v2	%x3 = load %x2		
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%v4 = icmp eq %v3, 0	%x4 = icmp eq %x3, 73	1
br %v4, L _{b3} , L _{b2}	br %x4, L _{b3} , L _{b2}	1

Merged Cost: 10 ✓

Main Differences Between CFM-CS and SEME-Fusion

Structural similarity

CFM-CS only works on branches with isomorphic SESE subregions.

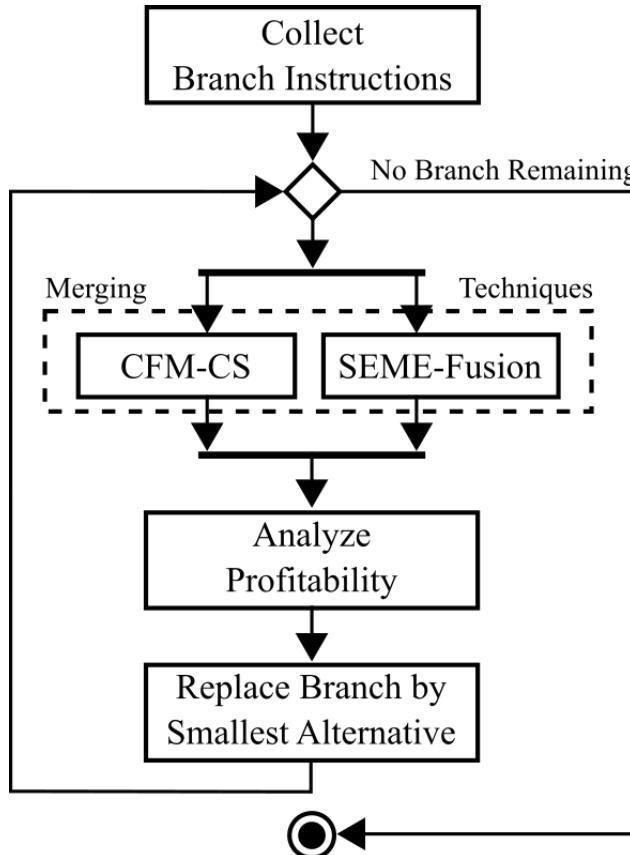
SEME-Fusion can merge any conditional branch - no structural similarity needed.

Block Paring

CFM-CS pairs the corresponding blocks in the isomorphic SESE subregions

SEME-Fusion pairs blocks based on fingerprint distance

HyBF: The Best of CFM-CS and SEME-Fusion



CFM-CS and SEME-Fusion compete,
the best result wins!

Evaluation Setup

Implemented in Clang/LLVM.

Benchmarks optimized with -Oz for code size.

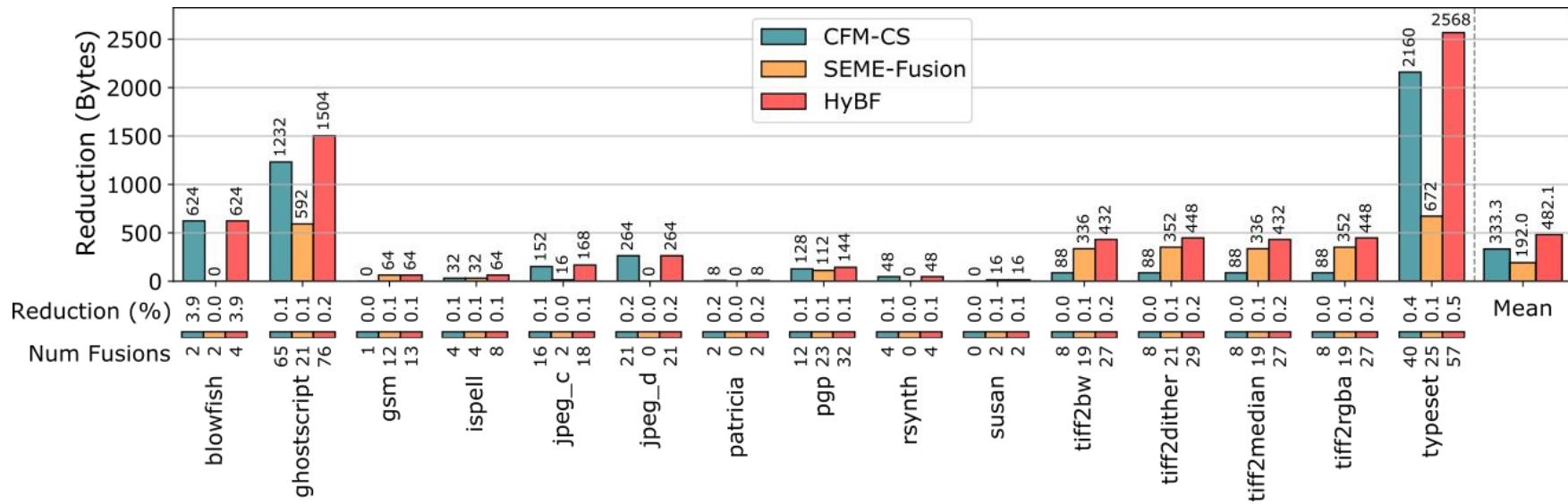
CFM-CS

SEME-Fusion

HyBF: Combines both CFM-CS and SEME-Fusion

Size Reduction

HyBF achieves the best reduction.



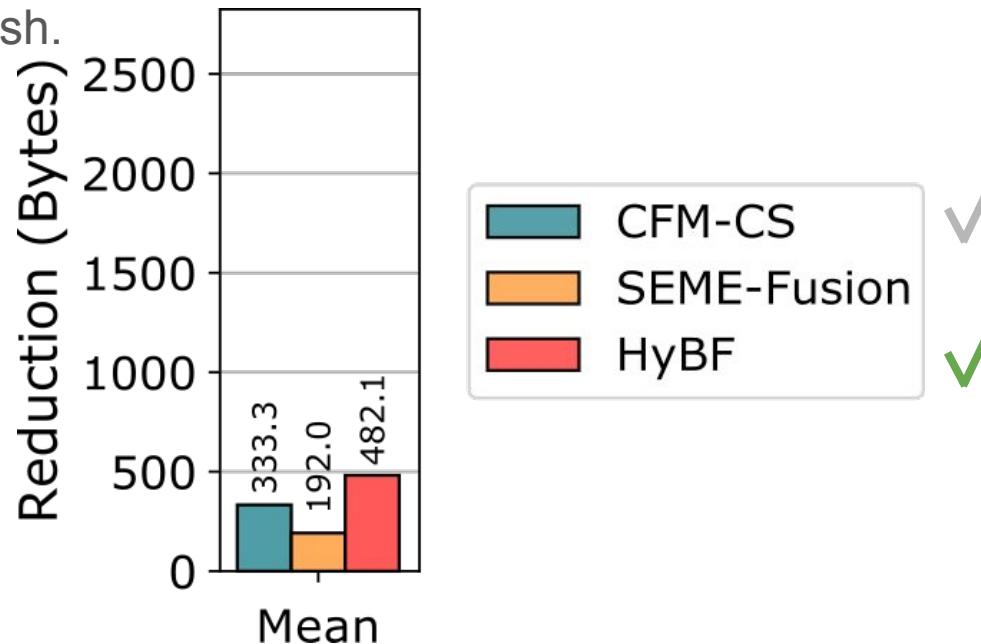
Mibench

Size Reduction

HyBF achieves the best reduction.

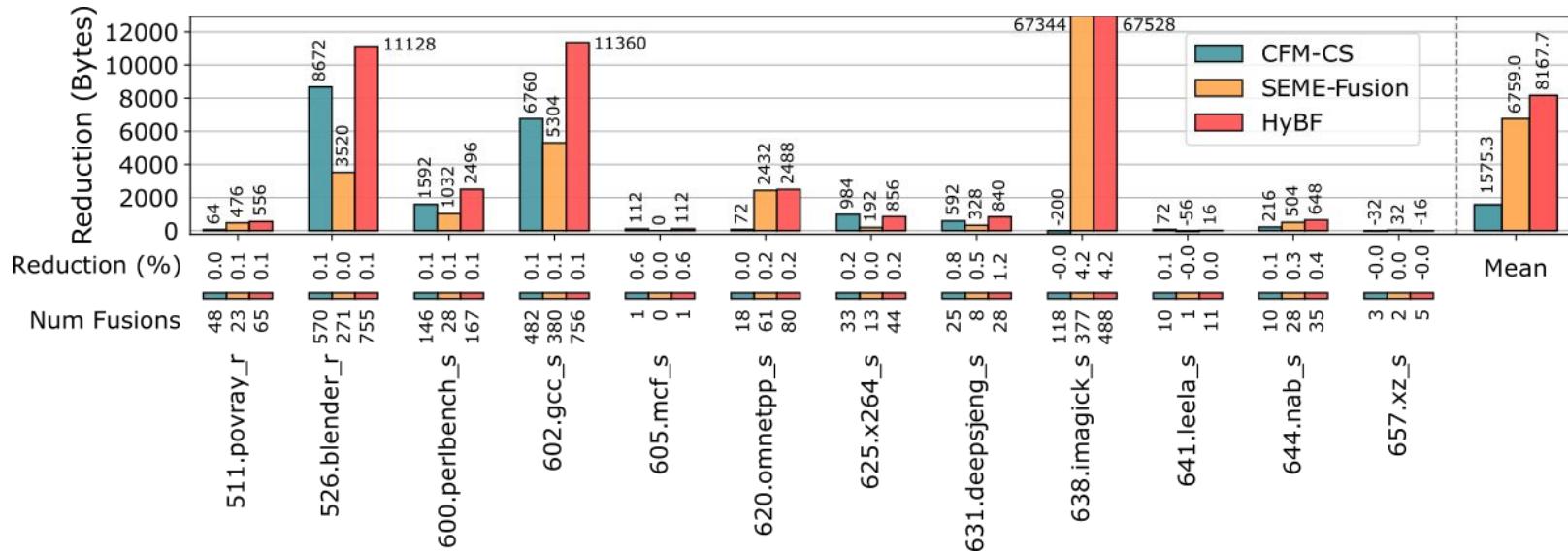
4% reduction in blowfish.

Mibench



Size Reduction

HyBF achieves the best reduction.



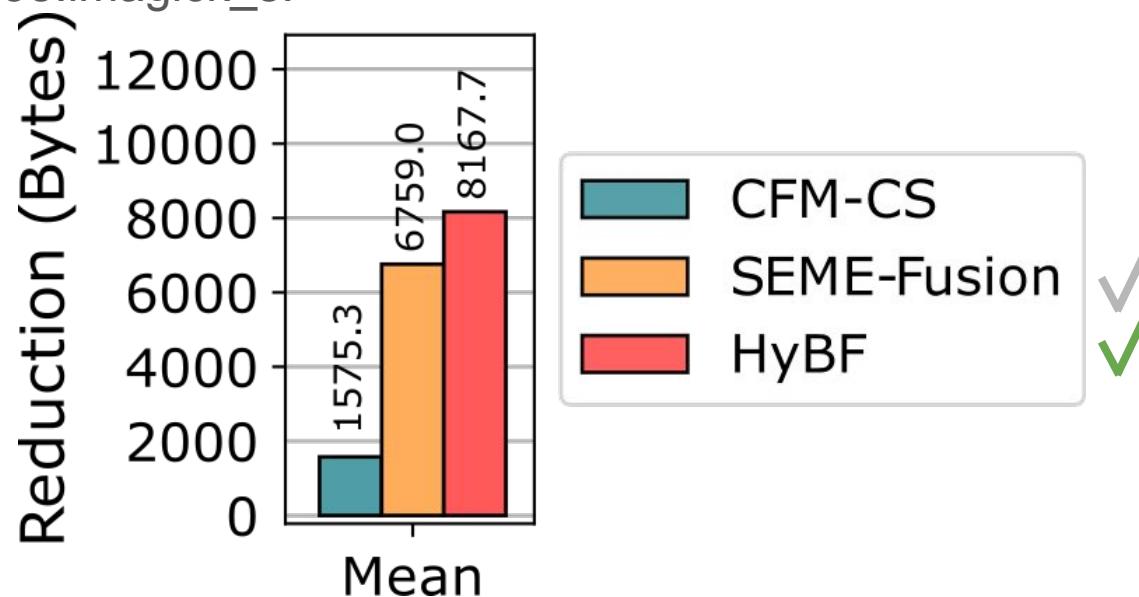
SPEC 2017

Size Reduction

HyBF achieves the best reduction.

4.2% reduction in 638.imagick_s.

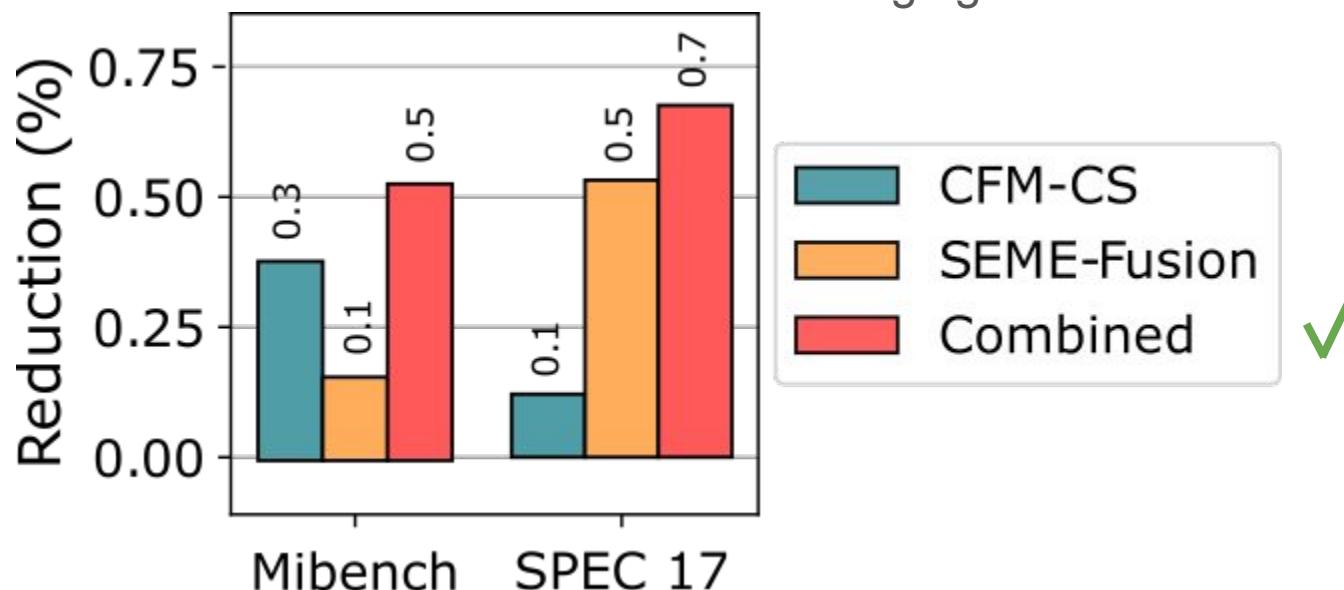
SPEC 2017



Size Reduction

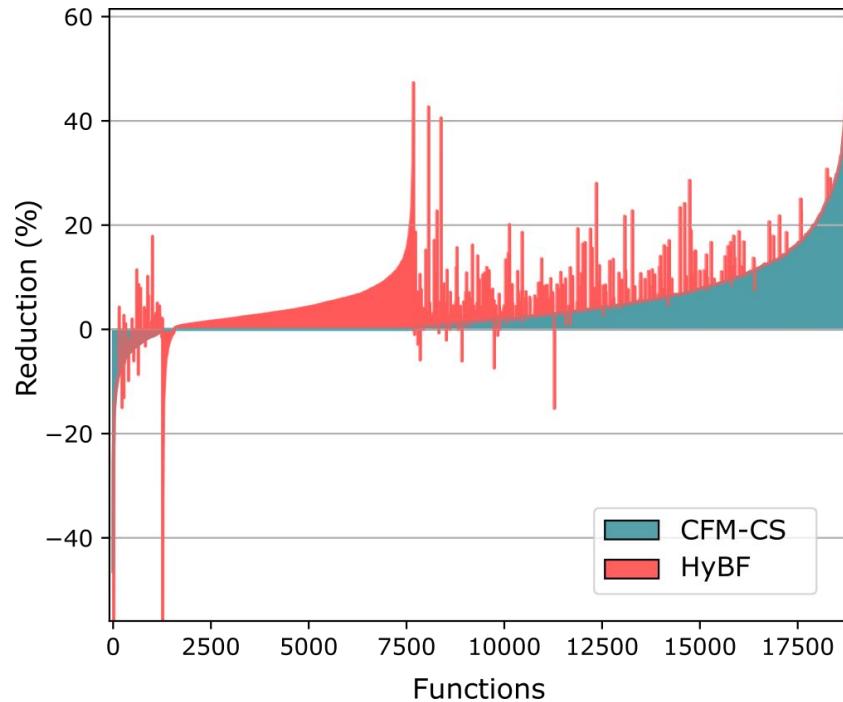
Branch Fusion is complementary to Function Merging:

All versions include the state-of-the-art function merging in LTO mode.



Size Reduction

HyBF tends to give better results than CFM-CS in AnghaBench functions.





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<https://github.com/charitha22/hybf-cc23-artifact>